# Local Area Networks (LANs) 

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## LAN Design Elements

- Four key elements of a LAN:
- Topology
- Transmission medium
- Wiring layout
- Medium access control


## Topologies

- Bus and Tree Topologies
- Used with multipoint medium
- Transmission propagates throughout medium and is heard by all stations
- Stations attach to the bus via a tap
- Full duplex connection between station and tap allows for transmission ("put on the bus") and reception ("take off the bus")
- Need to regulate transmission to avoid collisions and hogging
- Terminator absorbs frames at end of medium/cable
- Tree topology is a generalization of bus topology
- Headend connected to branching cables


## Bus and Tree Topologies


(b) Tree

## Frame Transmission on Bus Topology



Frame is not addressed to B; B ignores it


## Topologies

- Ring
- Point-to-point links between stations to form a loop Receive data on one link and retransmit on another
- Links unidirectional
- Stations attach to repeaters
- Data in frames
- Circulate past all stations
- Destination recognizes address and copies frame
- Frame circulates back to source where it is removed
- Media access control determines when a station can insert frame



## Frame Transmission on Ring Topology

1. $C$ transmits frame addressed to $A$

2. Frame is not addressed to $B ; B$ ignores it

3. A copies frame as it goes by


## Topologies

- Star
- Each station is directly connected to common central node
- Usually two point-to-point links between station and central node (one for transmit, one for receive)
- Two options:
- Hub: Central node broadcasts a frame to all stations (logically equivalent to a bus)
- Switch: Central node only sends frame to destination station (more efficient, but requires more complexity in central node)



## Choice of Topology and Medium

- Many factors: reliability, expandability, performance, building layout, medium available
- Some common cases (not always the case):
- Coaxial cable often used for bus topology
- Optical fibre for ring topology; usually the highest speed networks
- Twisted pair for star topologies; often well-suited for LANs in buildings


## IEEE 802 LAN Architecture

- IEEE 802 LAN/MAN standards committee has developed the majority of the LAN standards in use including:
- Ethernet, Fast Ethernet, Gigabit Ethernet, Token Ring, ...
- The 802 series of standards follow a common architecture
- Standardised only at Physical layer and Data Link layer
- Data Link layer is broken into two sub-layers: Medium Access Control (MAC) and Logical Link Control (LLC)
- Assumes any upper layers (e.g. any Network layer, including IP)
- Data Link layer must support multiple nodes accessing the link
- 802 can support many MAC/Physical protocols, and uses one common LLC protocol
- E.g. 802.2 defines the LLC; 802.3 defines MAC and Physical protocols for Ethernet; 802.11 defines the MAC and Physical protocols for wireless LAN


## IEEE 802 LAN Architecture

## OSI Reference Model



## Need for Multiple Access Techniques

- Currently, we have assumed data link layer connects two computers, that is, point-to-point topology
- Examples:
- one router has a optical fibre to another router
- a PC has a twisted pair cable to a router
- Two computer devices have dedicated link between them
- But in some cases have point to multipoint topology:
- Examples:
- An Ethernet with many computers connected to one cable
- Wireless systems: Wireless LAN, Mobile Phone, Satellite, ...
- Many computer devices (more than two) must share the link
- Multiple access techniques define rules for "sharing" the link
- That is, which devices can transmit so that they don't interfere with other transmissions


## Multiple Access: Where and How?

- Where is the control point?
- Centralised: a single node in the network (usually a special node) controls who transmits and how.
- Advantages are:
- Greater control over providing priority and guarantees
- Each node (except central node) is simple
- Avoids problems of distributing information between nodes
- Disadvantages are:
- Single point of failure; if central node fails, MAC won't work
- Central node may become a bottleneck, reduce performance
- Distributed: all nodes collectively make decisions on who transmits and how
- Advantages and disadvantages: opposite to centralised
- How to perform control?
- Synchronous (fixed): a specific capacity is dedicated (reserved) for a connection; similar to TDM, FDM, circuit switching
- Asynchronous (dynamic): capacity is allocated to nodes (connections) on demand
- Dynamic is preferable for LANs because difficult to predict traffic requirements of users
- Subdivided into: Round Robin; Reservation; Contention


## Round Robin MAC

- Each station in turn is given the opportunity to transmit
- E.g. Station 1 has an opportunity to transmit, then station 2 has an opportunity and then station 3 and so on
- If a station accepts to opportunity, it can transmit up to a maximum amount of time (or maximum amount of data)
- The station may reject the opportunity (e.g. if it has no data to transmit); then the next station is given the opportunity
- Efficient when:
- Many stations have data to transmit over a long period of time
- Inefficient of only some stations have data to transmit, because waste time switching between stations


## Reservation-based MAC

- Time is divided into slots, and stations reserve future slots for an extended or even indefinite period
- Similar to Time Division Multiplexing (TDM)
- E.g. If we have ten slots per second, then a station may reserve the second slot for the next 5 seconds
- Reservations may be centralised or distributed
- Suitable for stream traffic since data needs to be transmitted at regular, known intervals
- Voice/video calls, audio/video streaming


## Contention-based MAC

- No control over stations; each station contends (or "fights") for its chance to transmit
- E.g. random access: if a station has data to send, and no-one else is transmitting, then the station transmits; if someone else is transmitting, then wait a random period and try again
- Works well for bursty traffic (data arrives at unknown intervals, and sometimes in bursts)
- Simple to implement
- Distributed - each station follows its own rules
- Efficient for light/moderate loads (sometimes inefficient if heavy load)
- What is load? Amount of traffic.
- Although MAC protocols have been implemented using Round Robin, Reservation and Contention, the most common products today are roundrobin and contention


## LAN Elements

Bridges, Hubs, Switches

## LAN Elements: Bridges

- Almost every LAN needs to interconnect to other LANs and WANs. Two approaches:
- Bridges: for interconnecting LANs (most often using the same LAN protocol)
- Routers: for interconnecting LANs and WANs (often different protocols) - covered in later lectures
- A basic bridge interconnects two (or more) LANs using the same protocol
- Why not just make one large LAN (instead of connect two small LANs)?
- Reliability: a bridge can partition networks; a failure on one LAN will not affect the other LANs
- Performance: Generally LAN performance declines with length of wire and number of stations; partitioning can be used to increase performance
- Security: Traffic from one LAN (e.g. for finance department) can be kept separate from the other (e.g. student) LAN, although some monitored traffic can pass through the bridge
- Geography: it may be easier to connect a LAN in one building with a LAN in another building via microwave bridge (rather than coaxial cable to make one large LAN)


## Bridge Operation

## LAN A



## Bridge Operation

- A bridge connecting two LANs, $A$ and $B$, performs the functions:
- Read all frames transmitted on A, and accept those addressed to any station on B
- Using the MAC, retransmit each frame on B
- Do the same for B-to-A traffic
- Some design aspects of a bridge:
- Bridge does not modify the content or format of frames; nor does it add any more headers; it simply copies the received frame and retransmits (it can do this easily because same LAN protocol is used)
- Bridge contains buffer space to meet peak demand; sometimes frames may arrive much faster than the bridge can retransmit (hence the frames have to be buffered)
- Bridge must know something about addressing and routing - at minimum, it must know which addresses are on LAN A and which addresses on LAN B
- Bridge can interconnect more than two LANs (in which case each bridge must know more addressing/routing information)


## Connecting Multiple LANs

- Station 1 transmits a frame destined to station 6
- Bridges 101, 107 and 102 will receive on LAN A
- Only bridge 102 should retransmit
- Then bridge 105 and 106 will receive on LAN C
- Only bridge 105 should retransmit



## LAN Elements: Hubs

- In the star topology, a hub retransmits a received frame to all other stations
- This is effectively a bus: if two stations transmit at a time, then there will be a collision
- Often hubs are used in wiring closets in buildings to connect all computers on one floor; this suits the typical wiring of buildings today (wires from each office go to a central wiring closet)



## LAN Elements: Switches

- Switch used in same configuration as hub but only retransmits on a single line (not all)
- Also called "layer 2 switch" or "switching hub"
- More complex than hub
- Hub simply transmits the frame to everyone
- Switch needs to look at destination of frame, and transmit only to that station
- But switch is more common today because:
- Increase in performance in switched LAN (see next slide)
- Easy to upgrade from bus or hub to switch: all the stations use the same protocol/hardware, just connect the cables to switch
- Easy to scale network by adding more ports to switch
- Switches can perform functions of bridges as well
- Today, most devices sold are switches that include bridge functionality


## Switch Performance



## High Speed LAN Technologies

- Some of the early standards for LANs were:
- IEEE 802.3 bus- or star-based Ethernet
- IEEE 802.5 ring-based Token Ring
- Ethernet, and its many enhancements, has become the most dominant LAN technology today
- Ethernet (10Mb/s), Fast Ethernet (100Mb/s), Gigabit Ethernet (1Gb/s), 10GBe (10Gb/s), ...
- Alternatives today include Fibre Channel (up to 4Gb/s) and IEEE 802.11 Wireless LAN (54Mb/s and up)

|  | Fast Ethernet | Gigabit Ethernet | Fibre Channel | Wireless LAN |
| :--- | :---: | :---: | :---: | :---: |
| Data Rate | 100 Mbps | 1 Gbps, 10 Gbps | $100 \mathrm{Mbps}-3.2 \mathrm{Gbps}$ | $1 \mathrm{Mbps}-54 \mathrm{Mbps}$ |
| Transmission Media | UTP, STP, optical Fiber | UTP, shielded cable, <br> optical fiber | Optical fiber, coaxial <br> cable, STP | $2.4-\mathrm{GHz}, 5-\mathrm{GHz}$ <br> microwave |
| Access Method | CSMA/CD | Switched | Switched | CSMA/Polling |
| Supporting Standard | IEEE 802.3 | IEEE 802.3 | Fibre Channel <br> Association | IEEE 802.11 |

## Ethernet MAC

- IEEE 802.3 Ethernet uses a contention-based MAC protocol called:
- Carrier Sense Multiple Access (CSMA) with Collision Detection (CD)
- Aimed to ensure only one station transmits at a time, so to avoid collisions
- A collision is when two (or more) frames are sent on the medium at the same time - they interfere with each other, and neither are received, meaning costly (in terms of performance) retransmissions are needed
- With CSMA/CD the utilisation is less than 80\% for more than 5 nodes; hence throughput of $10 \mathrm{Mb} / \mathrm{s}$ Ethernet is typically between 5 and $8 \mathrm{Mb} / \mathrm{s}$
- This is shared between stations: with 10 stations, each station gets $800 \mathrm{~kb} / \mathrm{s}$
- However, Fast Ethernet uses a switched topology: each station has a dedicated link to the switch - there is no chance of collisions
- But the same IEEE 802.3 MAC frame format and protocol are used on Fast Ethernet so it is compatible with older 10Mb/s Ethernet
- LANs with mixture of $10 \mathrm{Mb} / \mathrm{s}$ Ethernet and $100 \mathrm{Mb} / \mathrm{s}$ Fast Ethernet devices can be deployed
- Throughput can approach 90\% with no collisions, hence each station gets $90 \mathrm{Mb} / \mathrm{s}$ with Fast Ethernet!


## IEEE 802.3 Specifications

|  | 10BASE5 | 10BASE2 | 10BASE-T | 10BASE-FP |
| :--- | :--- | :--- | :--- | :--- |
| Transmission <br> medium | Coaxial cable (50 <br> ohm) | Coaxial cable (50 <br> ohm) | Unshielded twisted <br> pair | 850-nm optical fiber <br> pair |
| Signaling <br> technique | Baseband <br> (Manchester) | Baseband <br> (Manchester) | Baseband <br> (Manchester) | Manchester/on-off |
| Topology | Bus | Bus | Star | Star |
| Maximum segment <br> length $(\mathbf{m})$ | 500 | 185 | 100 | 500 |
| Nodes per segment | 100 | 30 | - | 33 |
| Cable diameter <br> (mm) | 10 | 5 | 0.4 to 0.6 | $62.5 / 125 \mu \mathrm{~m}$ |


|  | 100BASE-TX |  | 100BASE-FX | 100BASE-T4 |
| :--- | :--- | :--- | :--- | :--- |
| Transmission <br> medium | 2 pair, STP | 2 pair, Category <br> 5 UTP | 2 optical fibers | 4 pair, Category <br> 3,4, or 5 UTP |
| Signaling <br> technique | MLT-3 | MLT-3 | 4B5B, NRZI | 8B6T, NRZ |
| Data rate | 100 Mbps | 100 Mbps | 100 Mbps | 100 Mbps |
| Maximum <br> segment length | 100 m | 100 m | 100 m | 100 m |
| Network span | 200 m | 200 m | 400 m | 200 m |

## Example Ethernet Configuration



## Ethernet Addressing

- Ethernet (IEEE 802.3) uses an IEEE standardised 48 bit address
- Known as EUI-48 or MAC-48, but often referred to in general terms as "MAC address", "Hardware address", "Network Interface Card (NIC) address", "Ethernet address", and others
- These addresses are also used by other Data Link layer protocols: Bluetooth, Wireless LAN (IEEE 802.11), ATM, Token Ring (IEEE 802.5), FDDI, Fibre Channel, ...
- IEEE 48-bit addresses
- Globally unique
- IEEE typically assign a unique first 24 bits to companies that manufacturer cards (e.g. Dell, 3Com, Intel). Each company then assigns a unique last 24 bits to each device they manufacture
- Individual users can "override" the assigned address (hence possibly that addresses are not globally unique) - used in security attacks
- Usually represented as hexadecimal, e.g. 00:17:31:5A:E5:89
- IEE 64-bit addresses
- Newer addresses used in Firewire, ZigBee/IEEE 802.15.4 and IPv6

