## CSS322 – Modes of Operation Notes

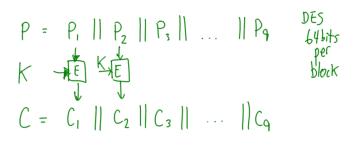


Figure 1: ECB Produces Repeating Ciphertext; Lecture 07

$$CBC : P_{1} = P_{2}$$

$$X_{1} = P_{1} \oplus IV$$

$$C_{1} = E(K, X_{1}) \qquad P_{1}$$

$$X_{2} = P_{1} \oplus C_{1}$$

$$C_{2} = E(K, X_{2}) \qquad P_{2}$$

$$Dces \quad C_{1} = C_{2} ?$$

$$C_{1} = E(K, P_{1} \oplus IV)$$

$$C_{2} = E(K, P_{1} \oplus C_{1})$$

$$if \quad IV = C_{1} \quad then \quad C_{1} = C_{2} \quad (BAD)$$

$$\uparrow$$

$$random$$

Figure 2: Conditions when CBC produces repeating ciphertext; Lecture 08

$$A \oplus B = C$$
,  $C \oplus B = A$ ,  $C \oplus A = B$   
 $(A \oplus B) \oplus B = A$ 



