

CSS322 – Public Key Cryptography

Notes

RSA Key Generation

$$p = 17 \quad q = 11$$

$$n = p \times q$$

$$= 17 \times 11$$

$$= 187$$

$$\phi(p) = p - 1$$

$$\phi(187) = \phi(17 \times 11)$$

$$= \phi(17) \times \phi(11)$$

$$= 16 \times 10$$

$$= 160$$

$$\gcd(e, 160) = 1 \quad 1 < e < 160$$

$$\cancel{2} \quad \textcircled{3} \quad \cancel{4} \quad \cancel{5} \quad \cancel{6} \quad \textcircled{7} \quad \textcircled{9} \quad \textcircled{11} \dots$$

$$e = 7$$

$$7 \times d \pmod{160} = 1$$

$$7 \times d = 161 \quad d = 23$$

$$PU_B = (e = 7, n = 187)$$

$$PR_B = (d = 23, n = 187)$$

Figure 1: RSA Key Generation; Lecture 10

$$\begin{array}{ccc}
 A & \xrightarrow{C=11} & B \\
 M=88 & & \\
 C = 88^7 \bmod 187 & & M' = 11^{23} \bmod 187 \\
 = 11 & & = 88 \quad \checkmark
 \end{array}$$

Figure 2: RSA Encryption and Decryption; Lecture 10

$$\begin{aligned}
 C &= M^e \bmod n & M &= C^d \bmod n \\
 M' &= C^d \bmod n \\
 &= (M^e \bmod n)^d \bmod n \\
 &= (M^{ed} \bmod n) \bmod n \\
 M' &= M^{ed} \bmod n \\
 a &= a^{ed} \bmod n \\
 a &= a^{\phi(n)+1} \bmod n \\
 ed &= \phi(n)+1 & 21 &= 20+1 \\
 e \times d \bmod \phi(n) &= 1 & 21 \bmod 20 &= 1 \\
 e : \text{RP}(e, \phi(n)) &\Rightarrow \text{gcd}(e, \phi(n))=1 \\
 \text{Calculate } d & & &
 \end{aligned}$$

Figure 3: Why RSA Successfully Decrypts; Lecture 11

$$\begin{array}{c}
 A \xrightarrow{C=11} B \\
 \downarrow \\
 \text{Attacker} \\
 \text{Know: } C=11, PU_B = (e=7, n=187) \\
 m = \underline{C}^{\textcircled{d}} \underline{\text{mod } n} \\
 ed \text{ mod } (\phi(n)) = 1 \\
 7 \times d \text{ mod } (\phi(187)) = 1 \\
 \phi(187) : \text{Factor } 187 \text{ into primes} \\
 RP(1, 187), RP(2, 187) \dots RP(186, 187)
 \end{array}$$

Figure 4: RSA Attack: Find totient; Lecture 11

$$\begin{array}{c}
 \text{Knows: } C=11, e=7, n=187 \\
 (m_1=17, C_1=85) \\
 \text{Find } d ?? \\
 \underline{m} = \underline{C}^{\textcircled{d}} \underline{\text{mod } n} \\
 d = d\log_{C,n}(m)
 \end{array}$$

Figure 5: RSA Attack: Find Discrete Log; Lecture 11

$$\begin{aligned}
 \text{PU} &= (e=7, n=299) \\
 \text{PR} &= (d=, n=299) \\
 ed &\equiv 1 \pmod{\phi(n)} \\
 7d \pmod{\phi(299)} &= 1
 \end{aligned}$$

$$\begin{aligned}
 n &= pq & n &= 299 \\
 \phi(n) &= \phi(pq) & &= pq \\
 &= \phi(p)\phi(q) & &= p-1 \\
 &= (p-1)(q-1) & &= q-1 \\
 \phi(299) &= 12 \times 22 & & \\
 &= 264 & &
 \end{aligned}$$

$$7x \pmod{264} = 1$$

$$7x = 265x$$

$$7x = (264 \times 2) + 1x$$

$$7x \underline{151} = (264 \times 4) + 1 \quad \checkmark$$

$$\underline{d} \quad \checkmark$$

Figure 6: Example of Breaking RSA Key Pair; Lecture 12