

Internet Protocols

ITS323: Introduction to Data Communications

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Internetworking Motivation and Requirements

The Internet Protocol

IP Addressing

LANs and WANs

LANs

- ▶ Different types: different topologies, different technologies, different purposes
- ▶ Many LANs operate at layers 1 and 2 (Physical and Data Link Layer) using switches and hubs
- ▶ Bridges can connect LANs of similar technologies together

WANs

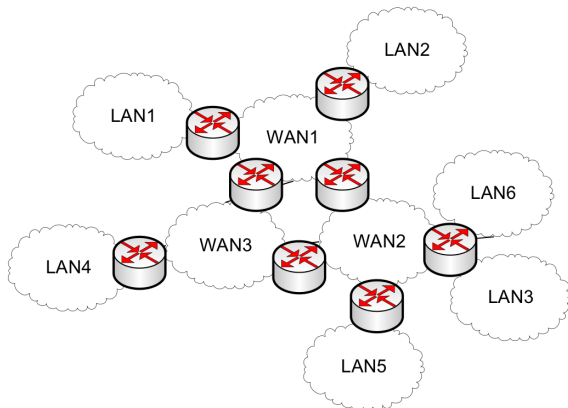
- ▶ Can interconnect LANs over a larger distance
- ▶ Point-to-point link (e.g. ADSL, PDH) or a network (e.g. ATM, SDH, telephone) using packet or circuit switching
- ▶ Device that interconnects the WAN to LAN must support both technologies
- ▶ WANs typically operate at Layers 1 and 2

Connect Multiple LANs and WANs

- ▶ Organisations have different requirements of their network, and therefore may choose different technologies for their LANs/WANs
- ▶ Aim: allow any computer to communicate with any other computer, independent of what LAN/WAN they are connected to
- ▶ **Internetworking** involves connecting the many different types of LANs/WANs together to achieve this aim

Internetworking with Routers

- ▶ Internetworking is performed using **routers**
- ▶ Routers connect two or more LANs or WANs together
- ▶ Routers are packet switches that operate at **network layer**



Terminology

- ▶ **Routers**: nodes that connect networks (LANs/WANs) together; operate at network layer
- ▶ **Subnetworks**: individual networks (LANs and WANs)
- ▶ **Internetworking**: connect two or more subnets together using routers
- ▶ An internetwork or an **internet**: the resulting network from internetworking
- ▶ **The Internet**: an internet that uses the Internet Protocol (IP) and used today to connect networks across the globe
- ▶ **Routing**: process of discovering a path from source to destination through a network
- ▶ **Forwarding**: process of sending data along a path through a network
- ▶ **Packet Switch**: a generic device that performs switching in a Packet Switching network. May operate at data link or network layer. A packet switch at network layer is called a router
- ▶ **Circuit Switch**: a generic device that performs circuit switching in a Circuit Switching network
- ▶ **Ethernet switch**: an IEEE 802.3 switch (either Ethernet, Fast Ethernet or Gigabit Ethernet). Operates at data link layer

Requirements of an Internetworking Protocol

- ▶ Provide link between subnetworks
- ▶ Provide for routing and delivery of data between processes on different subnets
- ▶ Provide service to keep track of use of networks and maintain status information
- ▶ Provide above services without requiring changes to the subnets. Accommodate differences between subnets, e.g.
 - ▶ Different addressing schemes
 - ▶ Different maximum packet size
 - ▶ Different timeouts
 - ▶ Error recovery
 - ▶ Status reporting
 - ▶ Routing techniques
 - ▶ Security
- ▶ The **Internet Protocol** meets some of these requirements. Others are left to **ICMP**, **TCP** and other protocols in the TCP/IP architecture

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The Internet Protocol

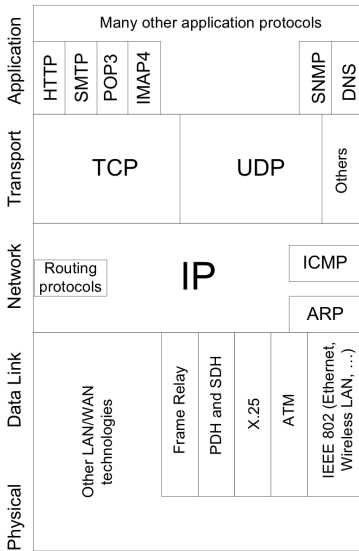
- ▶ IP is the internetworking protocol used in the Internet
 - ▶ We focus on IP version 4 (IPv4); IPv6 is available but not yet in widespread use
 - ▶ Other internetworking protocols: IPX, X.25, CLNP, SCCP
- ▶ Initially developed by US Department of Defence; now Internet Standard produced by IETF
- ▶ Features of IP:
 - ▶ Connectionless, network layer internetworking protocol using datagram packet switching
 - ▶ Provides data delivery, addressing, fragmentation and re-assembly
- ▶ Features IP does NOT provide:
 - ▶ Connection control, error control, flow control (TCP)
 - ▶ Status reporting (ICMP)
 - ▶ Priority, quality of service (DiffServ, IntServ)
 - ▶ Security (IPsec)

IP in the TCP/IP Stack

Motivation

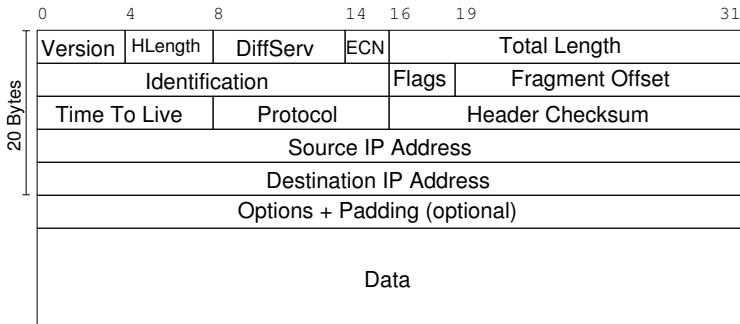
IP

IP Addresses



IP Datagram

- ▶ Variable length header and variable length data
- ▶ Header: 20 Bytes of required fields; optional fields may bring header size to 60 Bytes
- ▶ Data: length must be integer multiple of 8 bits; maximum size of header + data is 65,656 Bytes



IP Datagram Fields

- ▶ Version [4 bits]: version number of IP; current value is 4 (IPv4)
- ▶ Header Length [4 bits]: length of header, measured in 4 byte words
- ▶ DiffServ [6 bits]: Used for quality of service control
- ▶ ECN [2 bits]: Used for notifying nodes about congestion
- ▶ Total Length [16 bits]: total length of the datagram, including header, measured in bytes
- ▶ Identification: sequence number for datagram
- ▶ Flags: 2 bits are used for Fragmentation and Re-assembly, the third bit is not used
- ▶ Fragment Offset [13 bits]: See Fragmentation and Re-assembly
- ▶ Time To Live [8 bits]: datagram lifetime
- ▶ Protocol [8 bits]: indicates the next higher layer protocol
- ▶ Header Checksum [16 bits]: error-detecting code applied to header only; recomputed at each router
- ▶ Source Address [32 bits]: IP address of source host
- ▶ Destination Address [32 bits]: IP address of destination host
- ▶ Options: variable length fields to include options
- ▶ Padding: used to ensure datagram is multiple of 4 bytes in length
- ▶ Data: variable length of the data

Connectionless Internetworking with IP

Connection-oriented Internetworking

- ▶ Logical connection created between source and destination for data transfer
- ▶ All datagrams sent within connection are associated with each other
- ▶ Connection setup, data transfer, connection termination

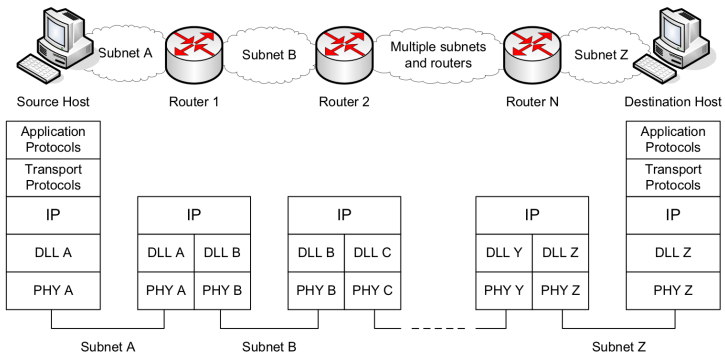
Connectionless Internetworking

- ▶ No connection between source and destination
- ▶ Datagrams are treated independently
- ▶ Advantages:
 - ▶ Flexible: can deal with different networks, requires little of subnets
 - ▶ Very small overhead if connectionless transport (e.g. UDP) is used

IP Hosts and Routers

- ▶ **Hosts** are the end-devices (stations)
 - ▶ Usually only use single network interface at a time
 - ▶ Hosts do not forward IP datagrams
 - ▶ Either source or destination
- ▶ **Routers** are the datagram packet switches
 - ▶ Routers have two or more interfaces (since they connect LANs/WANs together)
 - ▶ Routers forward datagrams
 - ▶ Routers can act as a source or destination of datagrams (however this is mainly for management purposes)
- ▶ **IP routing** is the process of discovering the best path between source and destination
- ▶ **IP forwarding** is the process of delivering an IP datagram from source to destination

IP Hosts and Routers



IP Routing

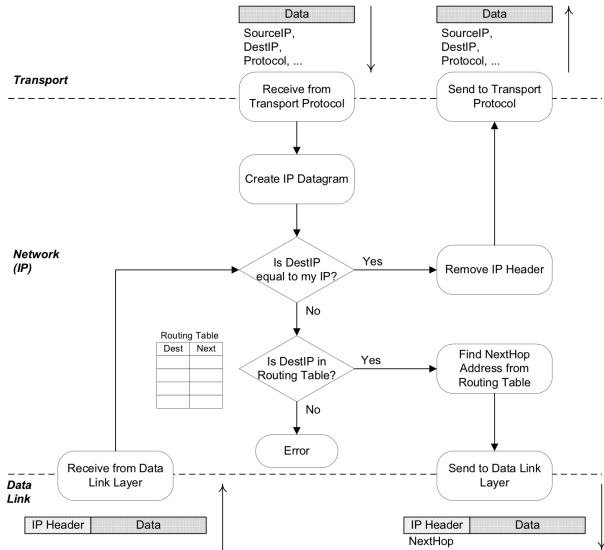
- ▶ No routing protocol is specified for IP
- ▶ Any of the available routing protocols can be used depending on the network topology and requirements of network administrator, e.g. RIP, EIGRP, OSPF, BGP, ...
- ▶ Each routing protocol creates and updates a routing table, which stores information to determine the path from source to destination
- ▶ IP uses the information in the routing tables to forward datagrams
- ▶ In order to make routing tables manageable, three strategies are used in the Internet:
 - ▶ Storing Next-Hop Routes
 - ▶ Network-specific Routing
 - ▶ Default Routes

IP Forwarding

Motivation

IP

IP Addresses



Fragmentation and Re-assembly

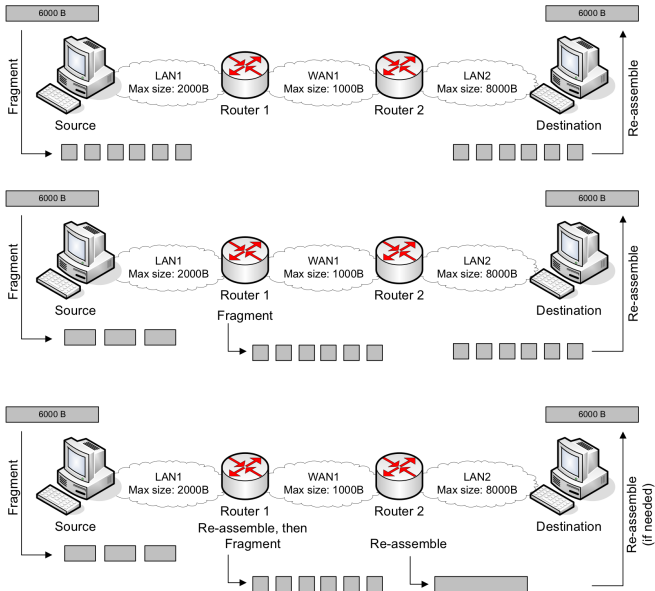
- ▶ Network layer may divide data from transport layer into multiple blocks (**fragmentation**)
- ▶ Data is **re-assembled** before being delivered to transport layer at destination
- ▶ Why fragmentation and re-assemble?
 - ▶ Subnets on path from source to destination may limit maximum size of frame
 - ▶ Error control may be more efficient with smaller packets
 - ▶ Smaller packets means smaller buffers needed at receivers
- ▶ Disadvantages of fragmentation and re-assembly:
 - ▶ Smaller packets means header contributes larger overhead
 - ▶ More packets means more time processing by routers, receiver

Fragmentation and Re-assembly

Motivation

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Fragmentation and Re-assembly

Three general options in internetworking:

1. Fragment only at source; re-assemble only at destination
2. Fragment at source and routers; re-assemble only at destination
3. Fragment at source and routers; re-assemble at routers and destination

IP uses option 2:

- ▶ No need for source to know maximum transmission units along path
- ▶ No need for routers to have large buffers for re-assembly
- ▶ No need for all fragments to pass through same router

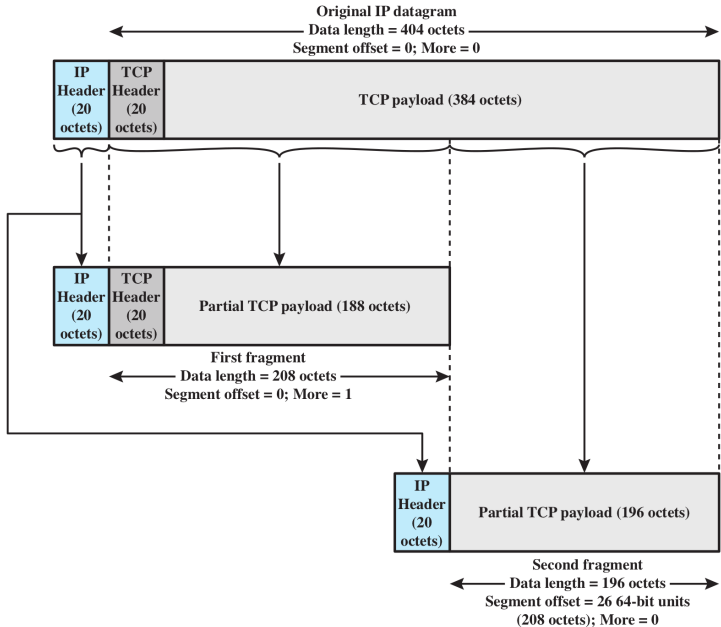
IP uses header fields to indicate if fragmentation has occurred and identify fragments

Fragmentation Example

Motivation

IP

IP Addresses



Datagram Lifetime in IP

- ▶ With adaptive/dynamic routing, it is possible for a routing loop; datagram sent forever
- ▶ Datagram marked with lifetime, when lifetime expires datagram is discarded
- ▶ IP uses a hop count:
 - ▶ Time To Live (TTL) field in header set to number of hops source allows the datagram to traverse (e.g. 64, 255)
 - ▶ Each router that processes datagram decrements the TTL field
 - ▶ If TTL is 0, datagram is discarded
- ▶ Simpler than using actual time, as would require synchronisation between clocks on devices

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IPv4 Addresses

- ▶ IPv4 addresses are 32 bits in length
- ▶ Split into **network** portion and **host** portion: first N bits identify a subnet in the Internet; last H bits identify an IP device (host/router) in that subnet
- ▶ All subnets in the Internet have unique network portion
- ▶ All IP devices in a subnet have same network portion, but unique host portions
- ▶ Where/how to split has changed over time: Classful, Subnet addressing, Classless addressing
- ▶ Focus on classless addressing
- ▶ Why split? Allows hierarchical addressing, makes routing in Internet scalable

Representing IPv4 Addresses

- ▶ Writing and remembering 32 bits is difficult for humans
- ▶ IP addresses usually written in **dotted decimal notation**
- ▶ Decimal number represents the bytes of the 32 bit address
- ▶ Decimal numbers are separated by dots

IP: 11000000111001000001000100111001

Classless IP Addressing

- ▶ **Subnet mask** or address mask identifies where the IP address is split between network and host portion
- ▶ Mask is 32 bits: a bit 1 indicates the corresponding bit in the IP address is the network portion; a bit 0 indicates the corresponding bit in the IP address is the host portion
- ▶ The mask can be given in dotted decimal form or a shortened form, which counts the number of bit 1's from left

IP: 10000010000100010010100110000001

Mask: 111111111111111111111000000000

Special Case IP Addresses

Selected IP addresses are used for special purposes; they cannot be used to identify a host

Network Address identifies a subnet in the internet; all bits in host portion are 0

Directed Broadcast Address identifies all hosts on a specific subnet; all bits in host portion are 1

Local Broadcast Address identifies all hosts on the current subnet; all bits are 1

Loopback Address identifies current host; first 8 bits are 01111111; also called localhost

Startup Source Address identifies host if currently it has no address; all bits are 0

Selected addresses reserved for private networks (e.g. not connected to Internet; behind NAT)

- ▶ 10.0.0.0—10.255.255.255
- ▶ 172.16.0.0—172.31.255.255
- ▶ 192.168.0.0—192.168.255.255

IP Addressing Example

View the IP address on your own computer.

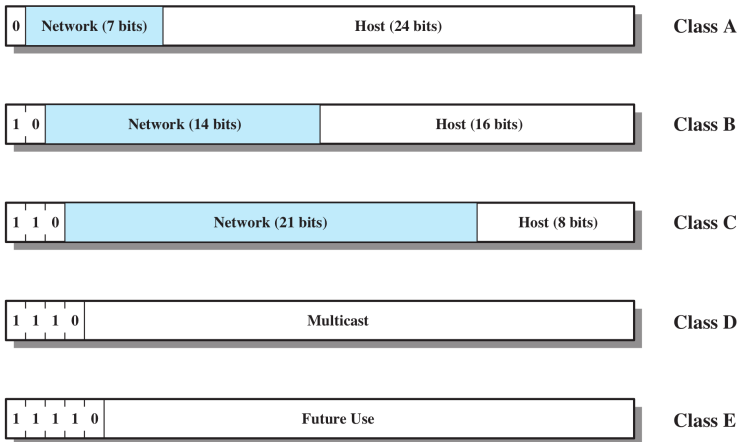
IP Addressing Example

My office computer has address 104.209.61.169/18. What is the network address and directed broadcast address for my network? How many IP devices can be attached to my network?

Classful IP Addressing

- ▶ Before “classless” addressing was introduced, the split was identified by one of 5 classes of addresses:
 - Class A first bit 0; network/host split after 8 bits (1.0.0.0—126.0.0.0)
 - Class B first two bits 10; split after 16 bits (128.0.0.0—191.255.0.0)
 - Class C first three bits 110; split after 24 bits (192.0.0.0—223.255.255.0)
 - Class D first four bits 1110; used only for multicast
 - Class E first five bits 11110; reserved for future use
- ▶ Subnet mask not needed; first bits of address determine the split
- ▶ Problem: only allow 3 different size networks (class A, B or C)

Classful IP Addressing



Obtaining an IP Address

- ▶ Internet Assigned Numbers Authority (IANA) manages the assignment of IP addresses
- ▶ IANA delegates IP network ranges to regional authorities (e.g. APNIC), delegated further to national registries (e.g. THNIC)
- ▶ Organisations obtain network addresses from national/local registries
- ▶ Organisations are free to assign addresses as they wish from assigned network address

Other Network Layer Functionality

- ▶ ICMP: error reporting, ping
- ▶ ARP: map IP addresses to Ethernet addresses
- ▶ IPv6
- ▶ Multicasting
- ▶ Quality of Service
- ▶ Mobility
- ▶ Security